Hierarchical Property Set Merging for SPARQL Query Optimization

Marios Meimaris, Athena Research Center, Greece

George Papastefanatos, Athena Research Center, Greece

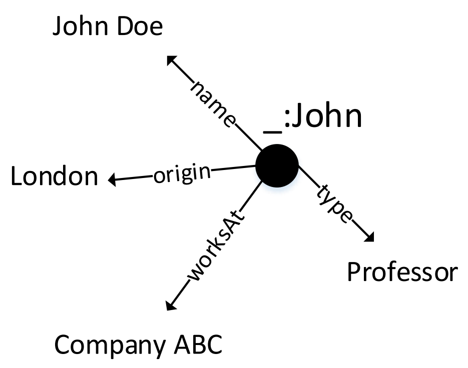
Panos Vassiliadis, University of Ioannina, Greece

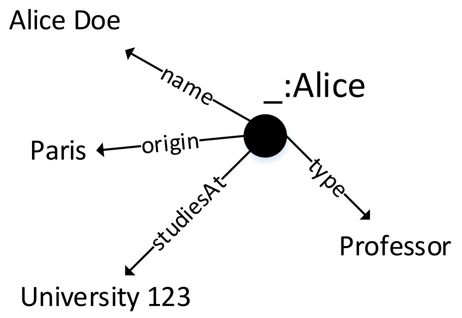
Preliminaries

- RDF (Resource Description Framework)
  - Abstract Data model for Linked Data
  - Based on *Triples*: Subject-Predicate-Object
  - RDF datasets are *Directed Labelled Graphs*
- Characteristic Set (CS)
  - A CS is a set of properties with the same subject as source node
  - An RDF dataset can be described as a set of unique CSs
  - Each CS is an *implicit resource type*

Preliminaries

- Use Characteristic Sets (CSs) and their links in order to store and index triples
- ***Characteristic Sets*** (Neumann & Moerkotte, ICDE 2011)
  - A *Characteristic Set* (CS) S_c of a node x is defined as the set of properties emitting from x (i.e., x as subject)



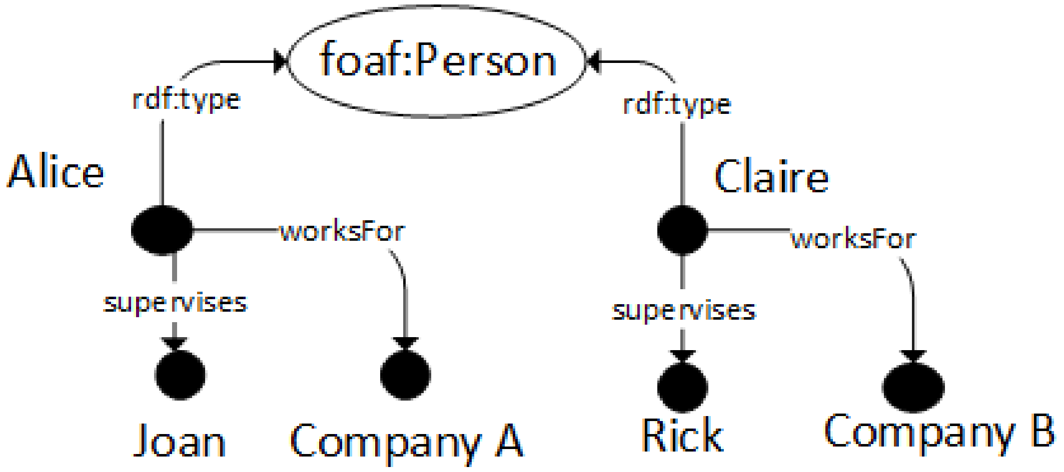


S_c(John) = {name, origin, worksAt, type}

S_c(Alice) = {name, origin, studiesAt, type}

Background

- Derive a relational representation of an RDF dataset
- Use CSs as tables and links between CSs as relationships
- CS properties → relation attributes

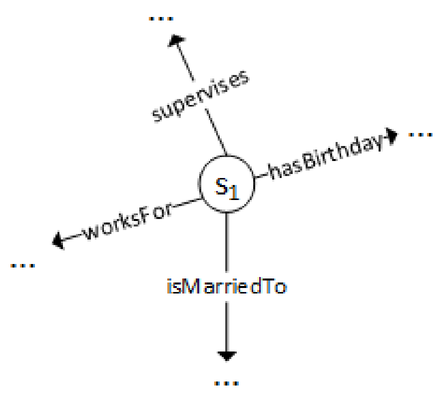


c_1= {rdf:type, worksFor, supervises}

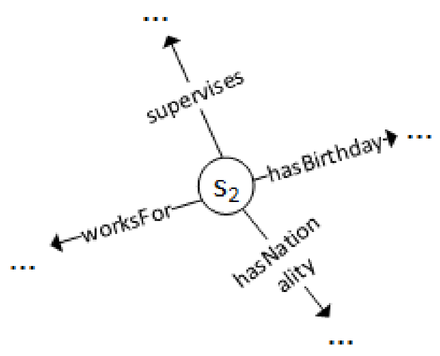
id	rdf:type	worksFor	supervises
Alice	foaf:Person	Company A	Joan
Claire	foaf:Person	Company B	Rick

Problem statement

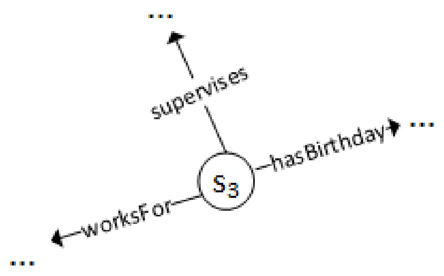
RDF structural looseness → multiple CSs → different representation strategies



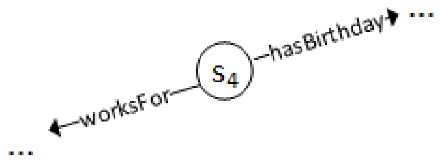
c_1 = {worksFor, supervises, hasBirthday, isMarriedTo}



c_2 = {worksFor, supervises, hasBirthday, hasNationality}

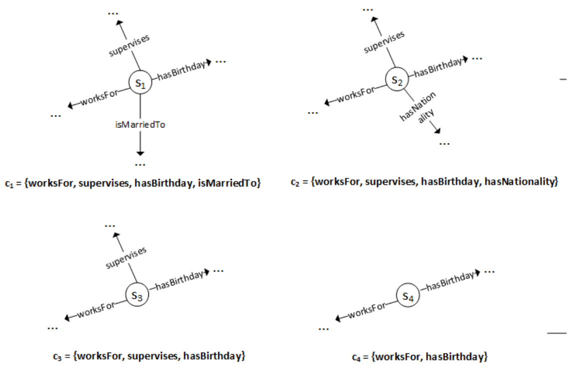


c_3 = {worksFor, supervises, hasBirthday}



c_4 = {worksFor, hasBirthday}

Trade-Off for creating a relational schema



A relational table for each different CS

id	supervises	worksFor	hasBirthday	isMarriedTo
s_1	...	...	...	...

id	supervises	worksFor	hasBirthday	hasNationality
s_2	...	...	...	...
...				

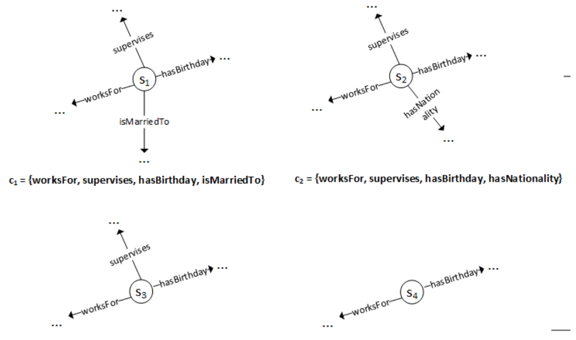
id	supervises	worksFor	hasBirthday
s_3	...	...	...
...			

id	worksFor	hasBirthday
s_4	...	...
...	...	...

A "universal" table for all CS’s

id	supervises	worksFor	hasBirthday	isMarriedTo	hasNationality
s_1	...	...	...	...	NULL
s_2	...	...	...	NULL	...
s_3	...	...	...	NULL	NULL
s_4	NULL	...	...	NULL	NULL
...	...	...	...	...	...

Trade-Off for creating a relational schema



c_1 = {worksFor, supervises, hasBirthday, isMarriedTo}
c_2 = {worksFor, supervises, hasBirthday, hasNationality}
c_3 = {worksFor, supervises, hasBirthday}
c_4 = {worksFor, hasBirthday}

A relational table for each different CS

id	supervises	worksFor	hasBirthday	isMarriedTo
s_1	...	...	...	...

id	supervises	worksFor	hasBirthday	hasNationality

- Space efficient
- Large numbers of relational tables with few tuples in
- Too many joins to answer queries

A "universal" table for all CS's

id	supervises	worksFor	hasBirthday	isMarriedTo	hasNationality
s_1	...	...	...	...	NULL

- Captures all CSs into a single table
- Too many NULL values
- Space inefficient

Trade-Off for answering a complex SPARQL query with many joins

```
SELECT ? x ?y ?z ?w
WHERE { ?x worksFor ?y .
        ?x supervises ?z .
        ?z hasBirthday '2011-02-24'.
        ?z isMarriedTo ?w.
        ?w hasNationality ‘GR’}
```

A relational table for each different CS

id	supervises	worksFor	hasBirthday	isMarriedTo
s_1	...	...	...	...

id	supervises	worksFor	hasBirthday	hasNationality

One self-join for each one of the *worksFor* , *supervises* and *isMarriedTo* query conditions

Additionally three joins between each CS table and all other CS tables in the database – i.e., 4 joins per table.

A "universal" table for all CS’s

id	supervises	worksFor	hasBirthday	isMarriedTo	hasNationality
s_1	...	...	...	...	NULL
s_2	...	...	...	NULL	...
s_3	...	...	...	NULL	NULL
s_4	NULL			NULL	NULL
...					

One self-join for each one of the *worksFor* , *supervises* and *isMarriedTo* query conditions

Problem to be solved

Context: Mapping heterogeneous RDF datasets to a relational schema with the aim to facilitate the processing of complex analytical SPARQL queries

Solution: automating the decision of which tables will be created for a set of CS, such that there are no overly empty tables and extremely large numbers of joins.

Observations

- Based on previous findings:
  - CS number is generally low but exhibits skewed distribution
    - E.g., many CSs with very few (<10) subjects
  - CS number affects number of joins
- Merging closely related CSs helps storage & querying
  - Less CSs means less joins
  - Less CSs means less I/O costs in disk-based systems
  - Compact schema easier to understand and maintain
- CSs are hierarchical, i.e., their property sets can be super/subsets of each other
- Challenge: *exploit the hierarchical structure in order to merge together closely related CSs*

Challenge

- Each CS defines a relational table $(s, p_1, p_2, \ldots, p_k)$
- Merging of CS tables results in NULL values for non-shared attributes
- Challenge: merge CSs and reduce NULL value effect

e.g.:

c_0 = {name, age}
c_1 = {name, age, marriedTo}
c_2 = {name, age, marriedTo, worksAt}

s	p_a	p_b
s_1	o_1	o_2
s_2	o_3	o_4

c_0

s	p_a	p_b	p_c
s_3	o_5	o_6	o_7
s_4	o_8	o_9	o_{10}

c_1

s	p_a	p_b	p_c	p_d
s_5	o_{11}	o_{12}	o_{13}	o_{14}
s_6	o_{15}	o_{16}	o_{17}	o_{18}

c_2

s	p_a	p_b	p_c	p_d
s_1	o_1	o_2	NULL	NULL
s_2	o_3	o_4	NULL	NULL
s_3	o_5	o_6	o_7	NULL
s_4	o_8	o_9	o_{10}	NULL
s_5	o_{11}	o_{12}	o_{13}	o_{14}
s_6	o_{15}	o_{16}	o_{17}	o_{18}

c_{merged}

Approach

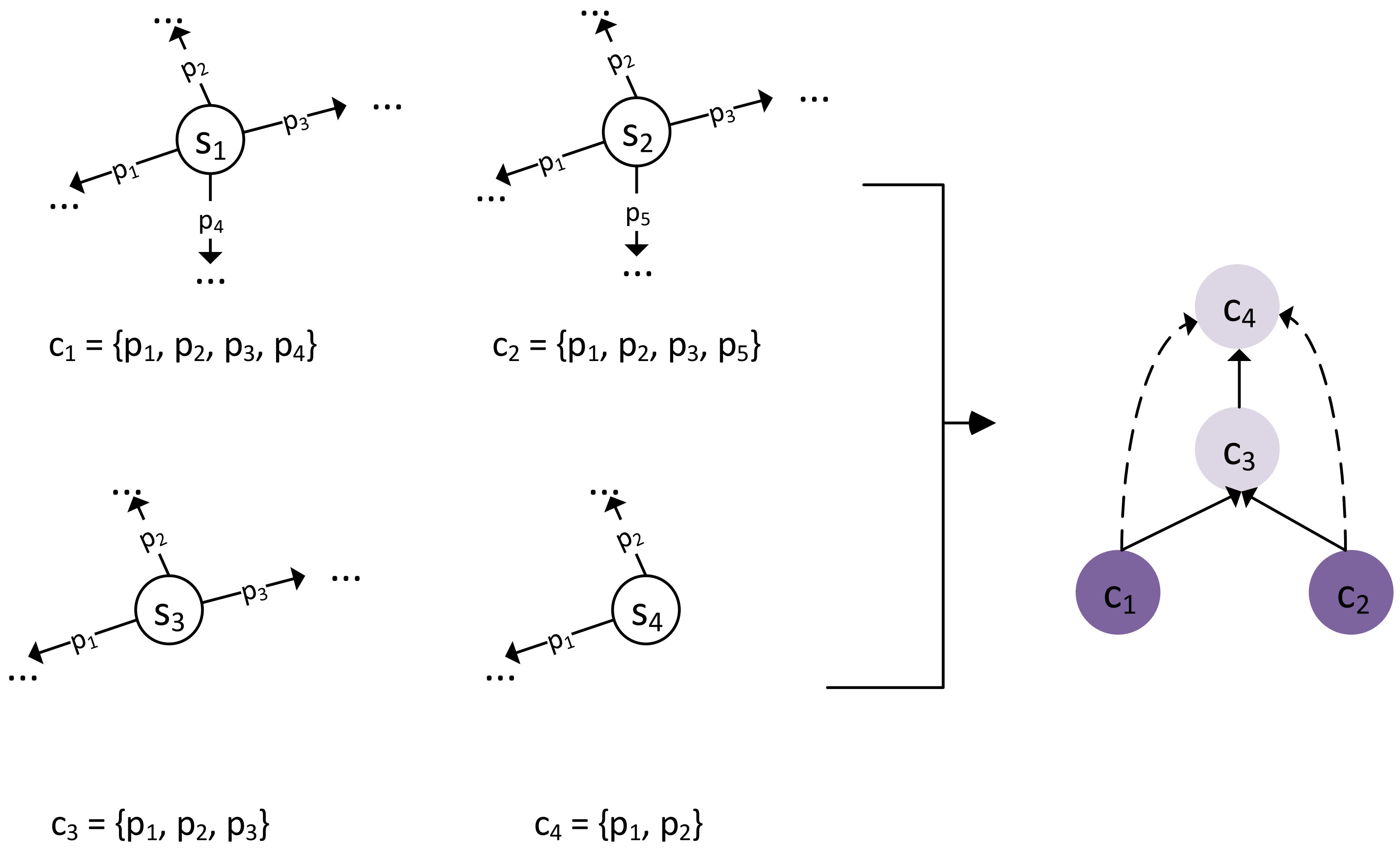
- Use a dense child table and merge its parents into it
  - **Why dense? ->** # of NULLs is proportional to # of records of table to be merged
  - **Why child? ->** more specialized, thus will contain columns of parents
- Identify dense CSs
    - if $|c_i| > m \times |c_{max}|$ parameter => c_i is dense
  - Every resulting (merged) table will contain **exactly one dense node** (and several non-dense)
- Find optimal merging of ancestors to dense child CSs

e.g.
c_1: {name, age, address}, c_2: {name, age}: c_1 child of c_2

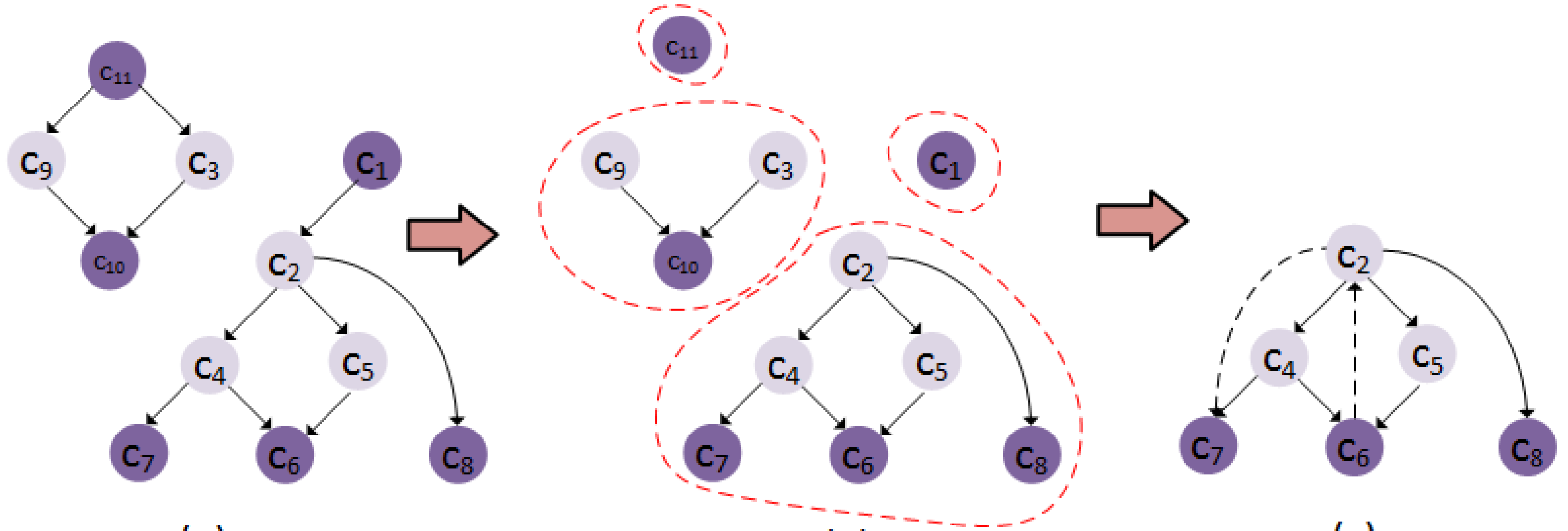
hier_merge(c_1, c_2) = c_{12}

c_{12}: {name, age, address)

CS Graph Example



Approach - Example



(a) (b) (c)

c_2	c_4	c_5
c_7	c_7	c_6
c_7	c_6	c_6
c_6	c_7	c_6
c_6	c_6	c_6
c_8	c_7	c_6
c_8	c_6	c_6

(d)

Approach – Loading and Merging

- Finding the optimal solution is equivalent to enumerating all possible sub-graphs -> *exponential*
- Greedy approximation
  - At each step, merge parent CS and dense child CS that minimize objective cost function
  - Cost function minimizes the number of NULL values introduced by the merge
- Tuning of *m* parameter

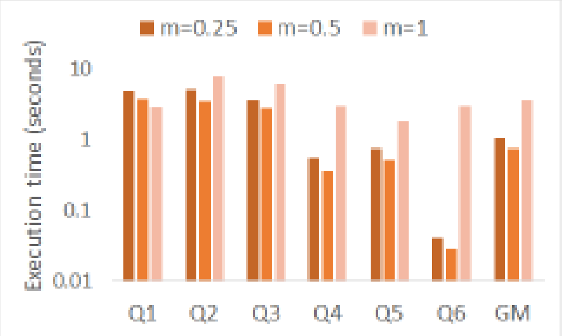
Approach – Querying

- Parse incoming SPARQL queries
  - Identify query CSs that match merged CSs in the dataset
  - Rewrite query as an SQL statement with UNIONs between matched CSs
  - In case of SO/OS joins, prune off CSs that are not linked
- Pass final query to relational optimizer
- Build and output results

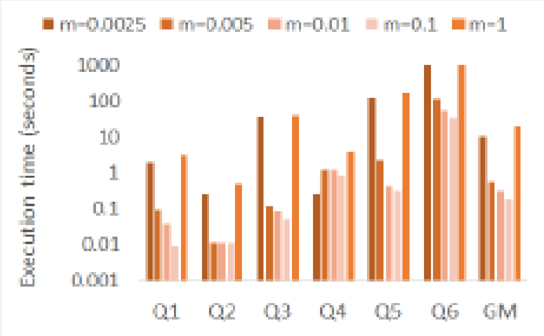
Implementation & Evaluation (Loading)

Dataset	Size (MB)	Time	# Tables (CSs)	# of ECSs	Dense CS Coverage
Reactome Simple	781	3min	112	346	100%
Reactome (m=0.05)	675	4min	35	252	97%
Reactome (m=0.25)	865	4min	14	73	77%
Geonames Simple	4991	69min	851	12136	100%
Geonames (m=0.0025)	4999	70min	82	2455	97%
Geonames (m=0.05)	5093	91min	19	76	87%
Geonames (m=0.1)	5104	92min	6	28	83%
LUBM Simple	591	3min	14	68	100%
LUBM (m=0.25)	610	3min	6	21	90%
LUBM (m=0.5)	620	3min	3	6	58%
WatDiv Simple	4910	97min	5667	802	100%
WatDiv (m=0.01)	5094	75min	67	99	77%
WatDiv (m=0.1)	5250	75min	25	23	63%
WatDiv (m=0.5)	5250	77min	16	19	55%

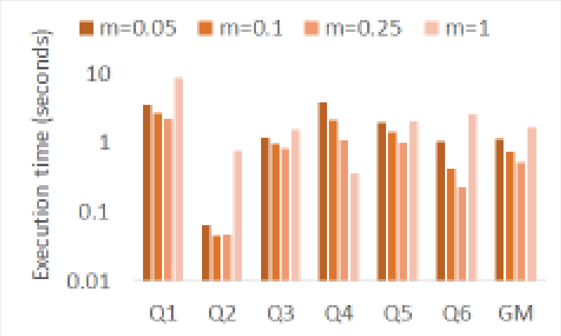
Implementation & Evaluation (Querying)



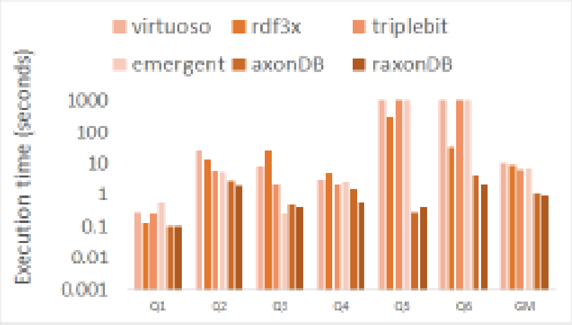
(a) Execution time (seconds) for LUBM



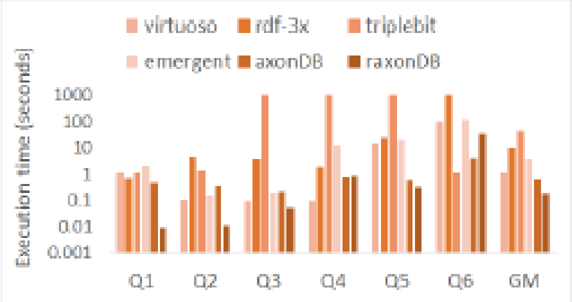
(b) Execution time (seconds) for Geonames



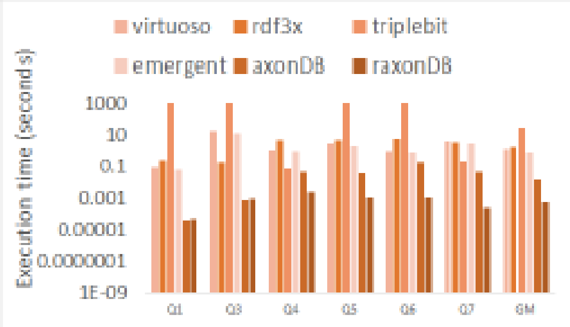
(c) Execution time (seconds) for Reactome



(a) Execution time (seconds) for LUBM2000



(b) Execution time (seconds) for Geonames



(c) Execution time (seconds) for Reactome

Future Work

- Distributed version of *raxonDB*
  - CS-based partitioning scheme
  - Distributed query processing
- Refined cost function
- Different ways of defining density

Thank you

{m.meimaris, gpapas}@athenarc.gr, pvassil@cs.uoi.gr

https://github.com/mmeimaris/raxonDB

https://visualfacts.imsi.athenarc.gr/

H.F.R.I.
Hellenic Foundation for Research & Innovation

This research is funded by the project VisualFacts (#1614) - 1st Call of the Hellenic Foundation for Research and Innovation Research Projects for the support of post-doctoral researchers.